

INVESTICE DO ROZVOJE VZDĚLÁVÁNÍ

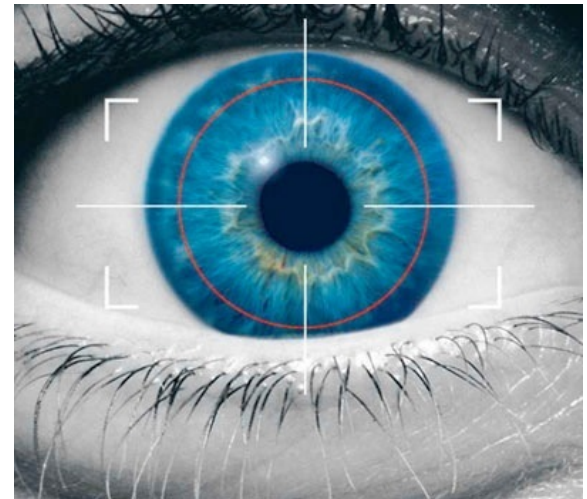
# Biometric Cryptography - Mobile Application Viewpoint

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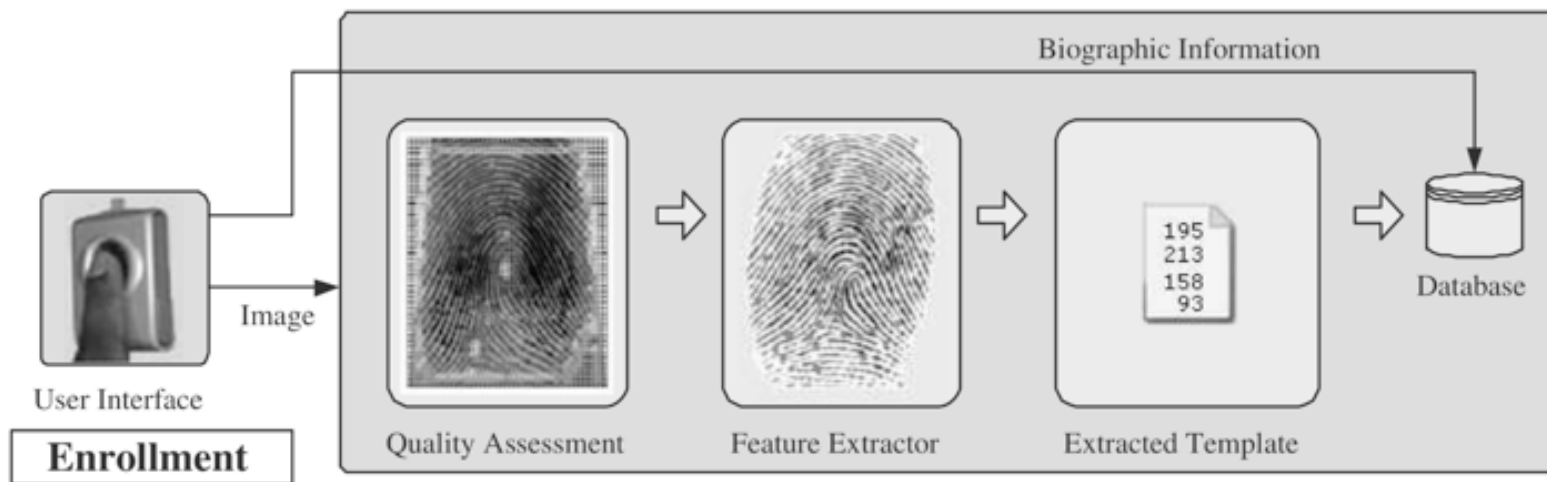


# Biometric Identification/Verification

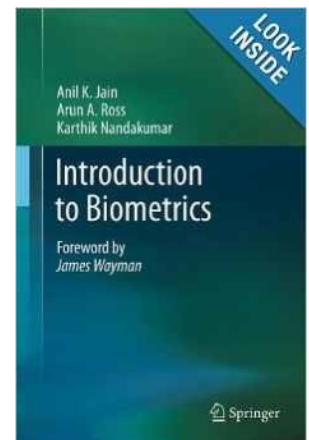
...automated establishment of the human identity based on their physical or behavioral characteristics.



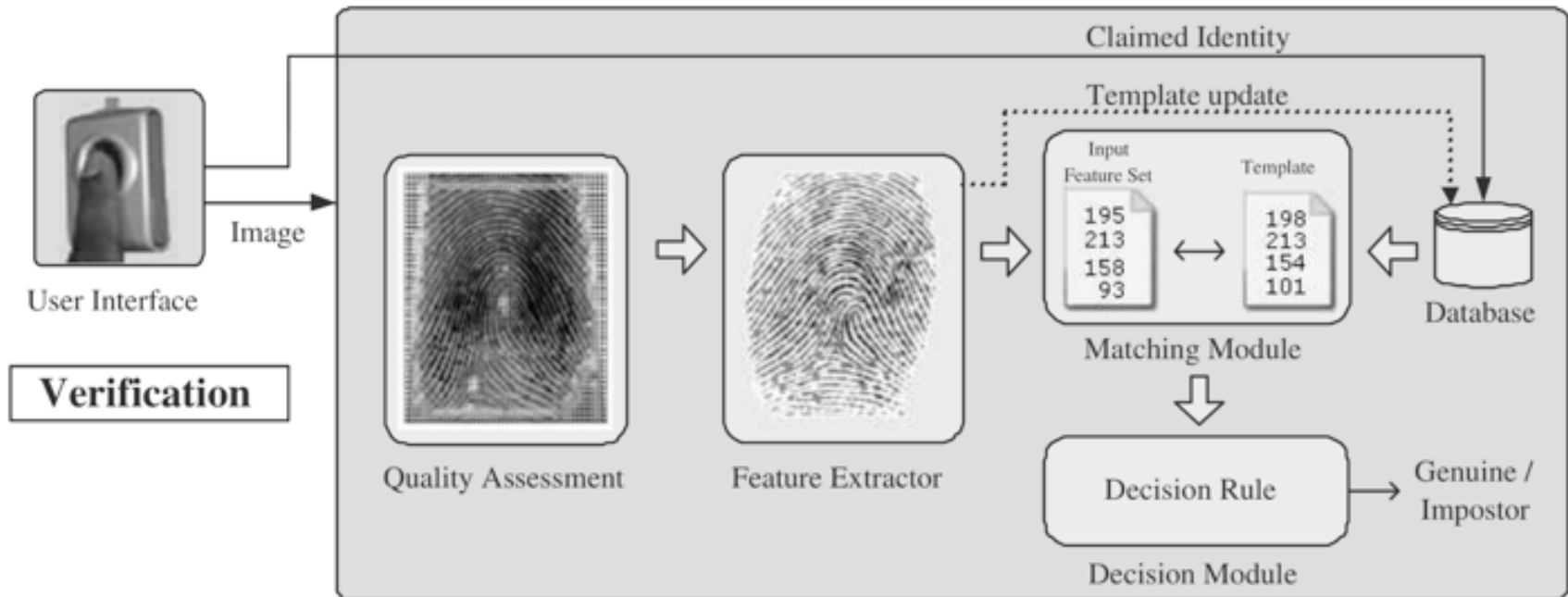
# [ Enrolment Phase ]



Jain, Ross, Nandakumar: *Introduction to Biometrics*, Springer, 2011

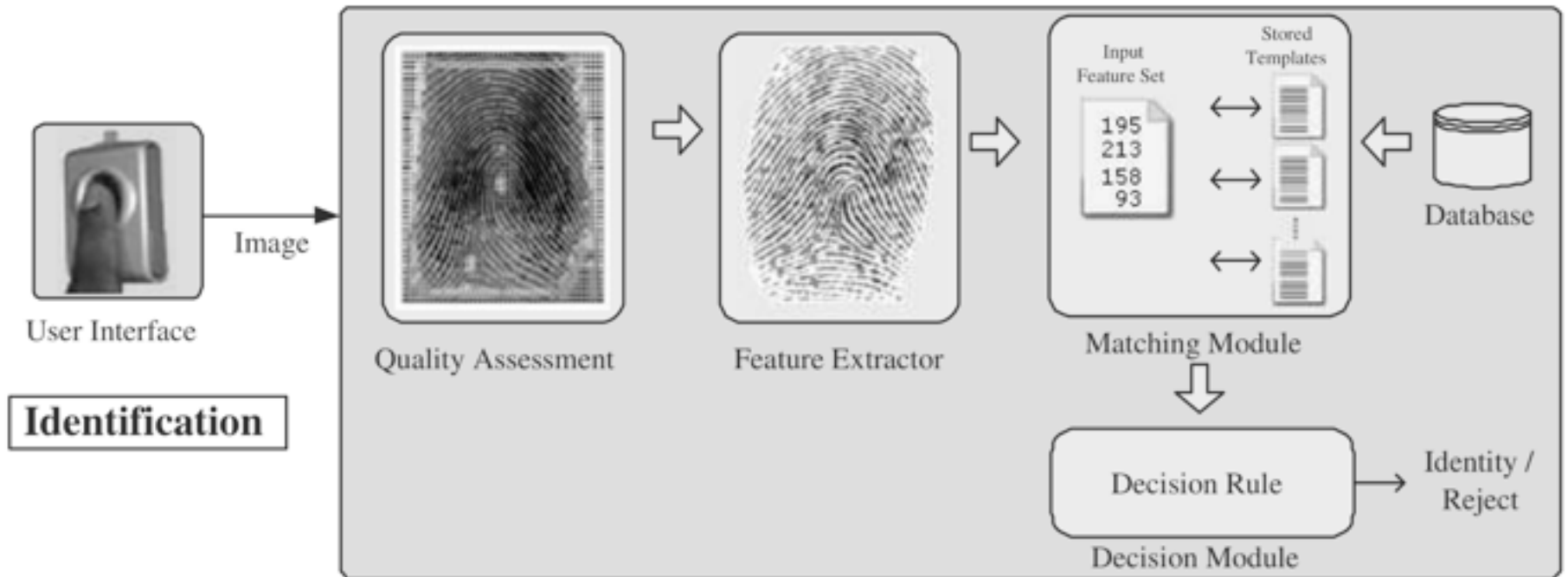


# Verification (1 : 1)



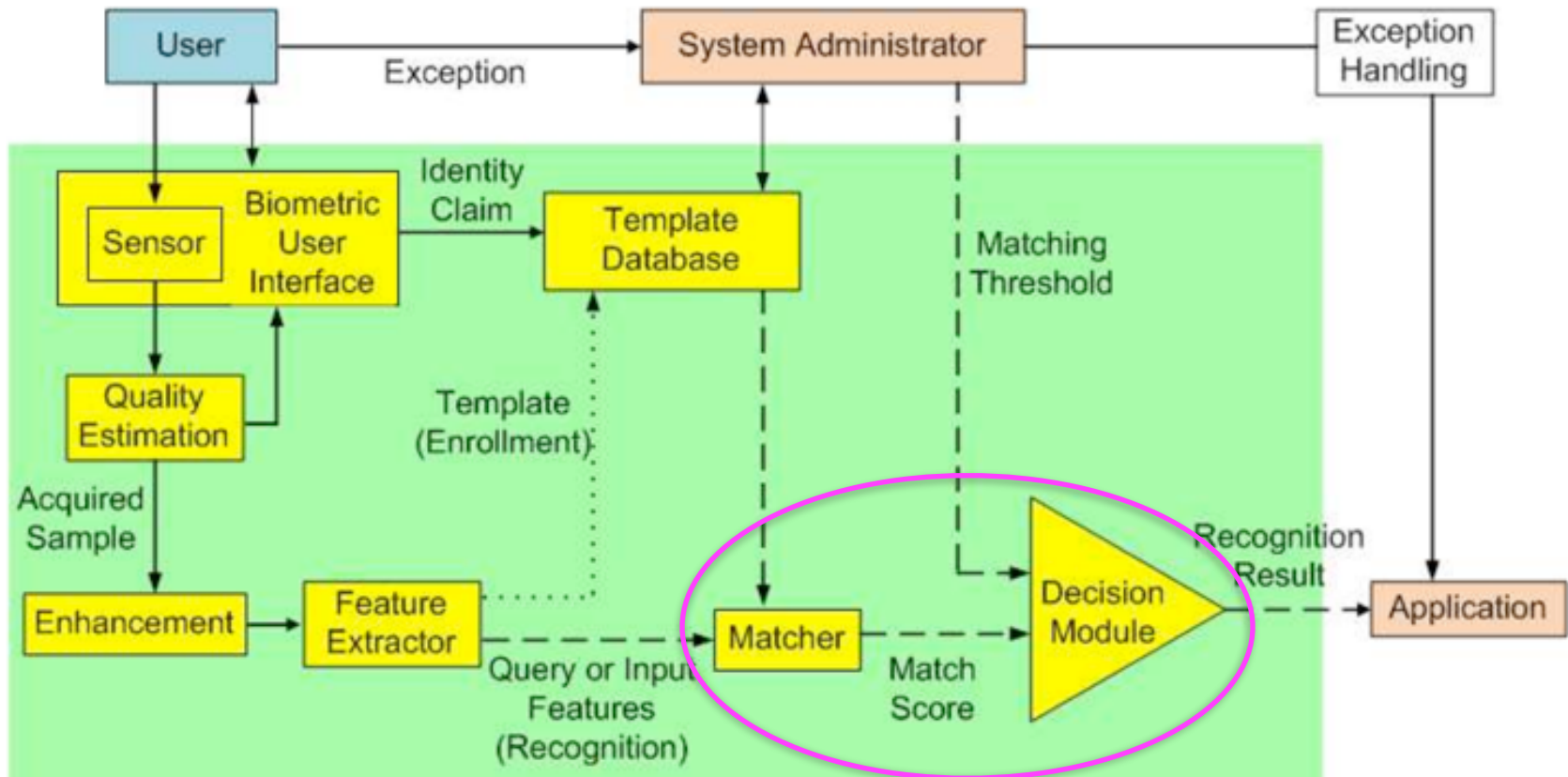
Jain, Ross, Nandakumar: *Introduction to Biometrics*, Springer, 2011

# [ Identification (1 : N) ]



Jain, Ross, Nandakumar: *Introduction to Biometrics*, Springer, 2011

# Biometric System Topology

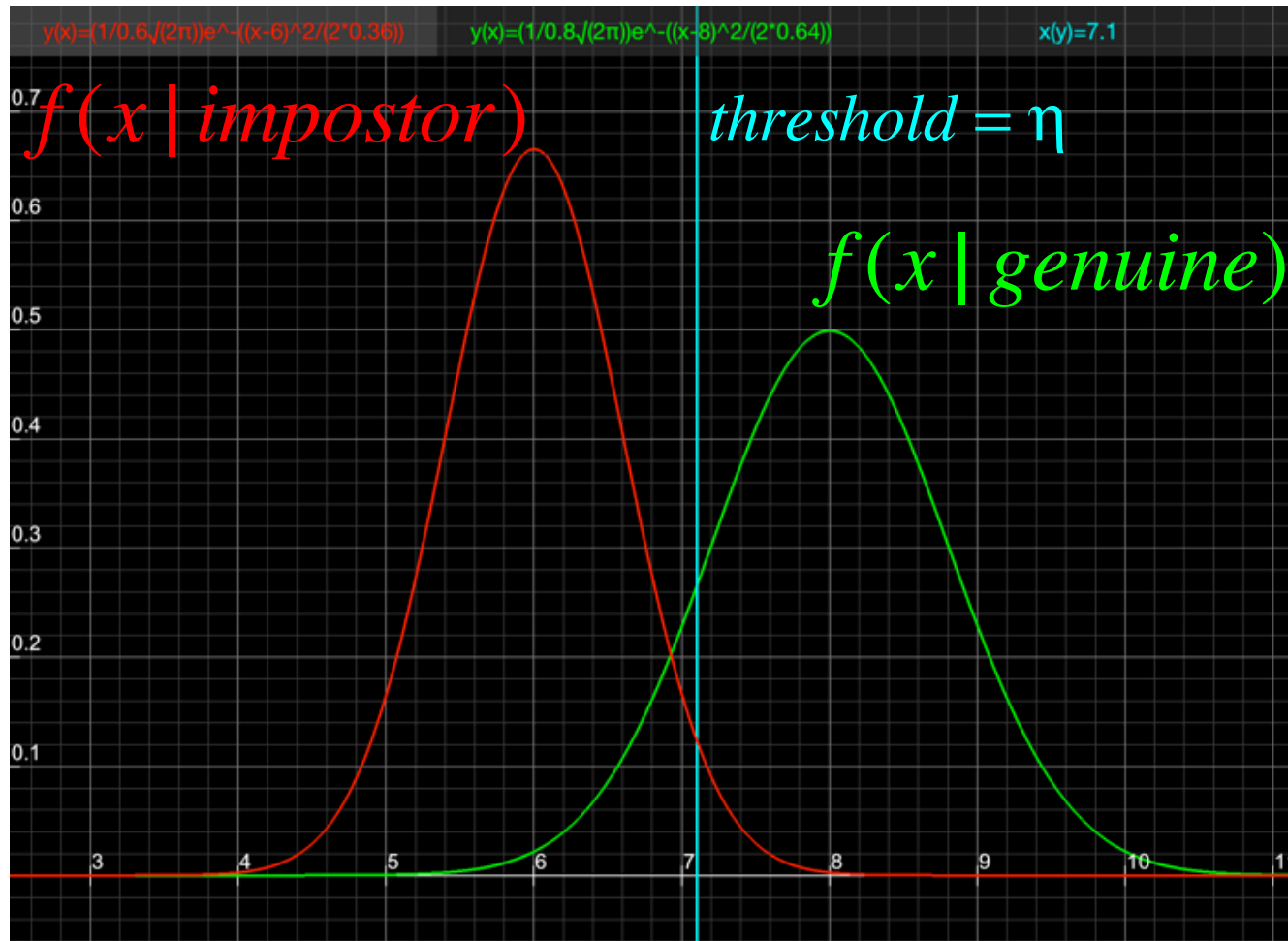


Jain, Ross, Nandakumar: *Introduction to Biometrics*, Springer, 2011

# [ Match Score ]

- It would be nice if we had simple **true/false** result.
  - As in conventional crypto.
  - But we cannot...
- All we have is a random variable  $X$  that follows two conditional distributions.
  - $f(x \mid \text{impostor})$
  - $f(x \mid \text{genuine})$

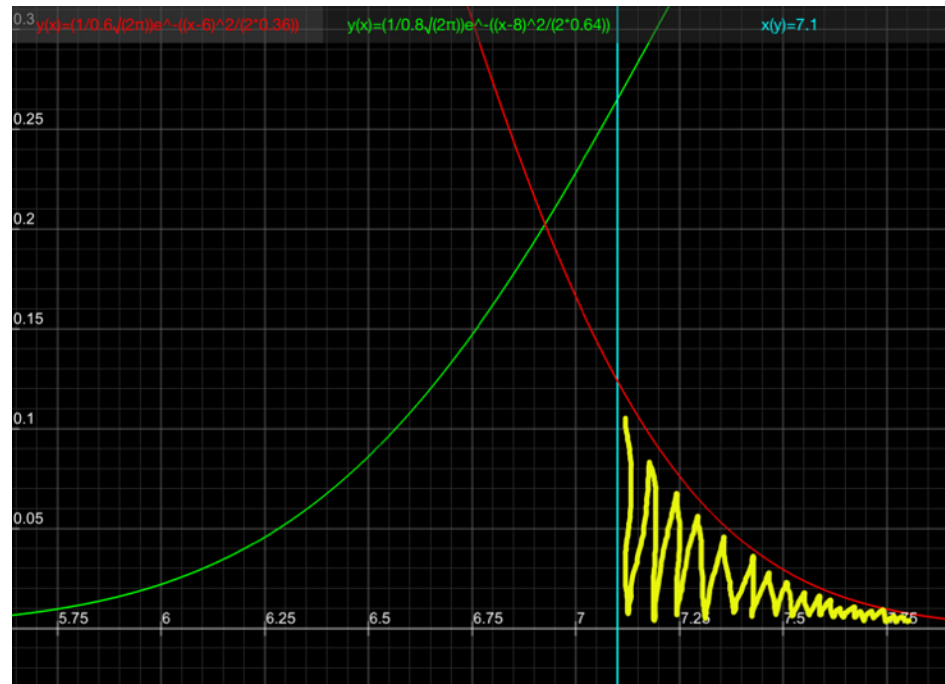
# [ Match Score Evaluation ]





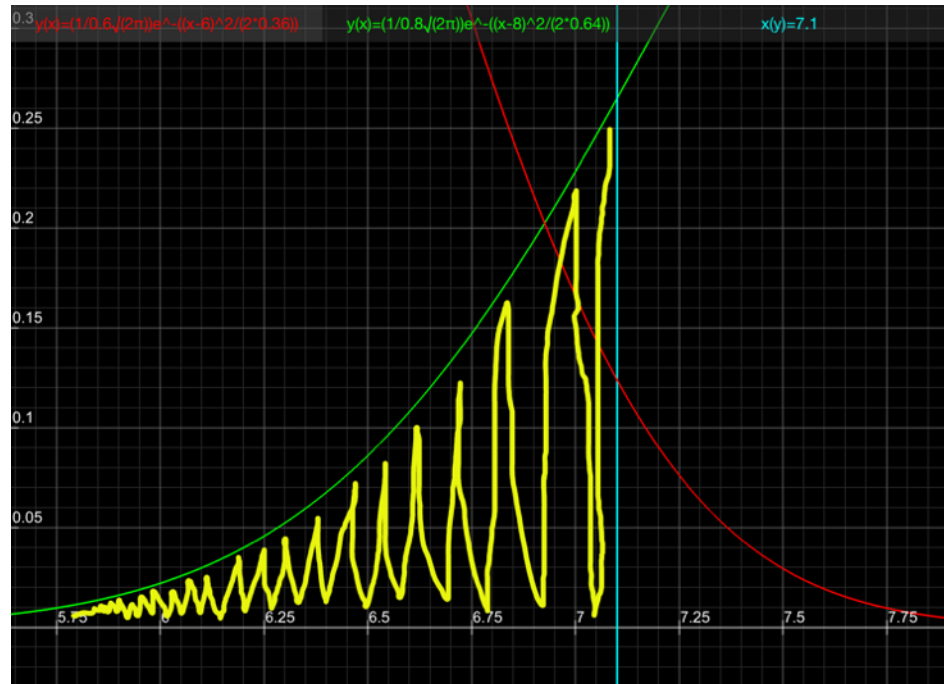
# [ False Acceptance Rate ]

$$FAR = \int_{\eta}^{\infty} f(x | impostor) dx$$

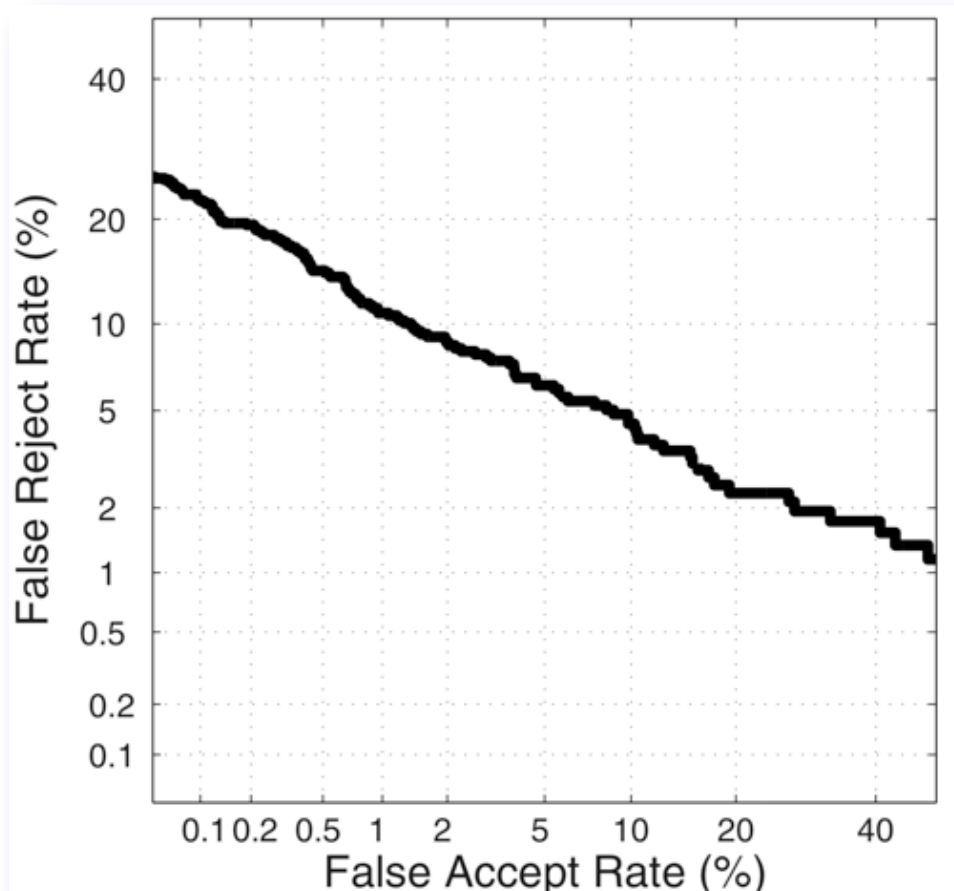


# [ False Rejection Rate ]

$$FRR = \int_{-\infty}^{\eta} f(x | \text{genuine}) dx$$



# [ Real DET Curve ]



## Detection Error Tradeoff

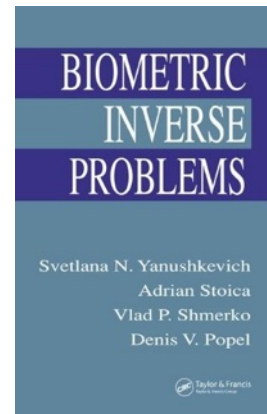
Jain, Ross, Nandakumar,  
Springer 2011

# [ Contrasting Design Approach ]

- Classic cryptography
  - infeasible mathematical problems
- Quantum cryptography
  - intractable physical problems
- Biometric identification
  - statistical signal analysis and pattern recognition
  - intractability is usually *not* the prime concern
  - we hope the Mother Nature complexity *somehow* guarantees the security

# [ BIO Brute Force Attack ]

- Randomly generate plausible circa 1/FAR samples and put them to the test.
  - Also termed “Zero-Effort”, denoting that the attacker makes no special effort to imitate the original person characteristic.
- Synthetic samples generation is quite feasible today.



# [ Cryptanalysis-Like Attacks ]

- Usually a variant of “Hill-Climbing” denoting the attacker iteratively improves the BIO sample data based on:
  - scoring feedback (*side channels*)
  - stolen template (*pre-image attacks*)
  - independent template trained from intercepted BIO samples (*correlation attacks*)
  - known scoring anomaly (*differential analysis. etc.*)
  - implementation faults (*general hacking*)

# [ Spoofing ]

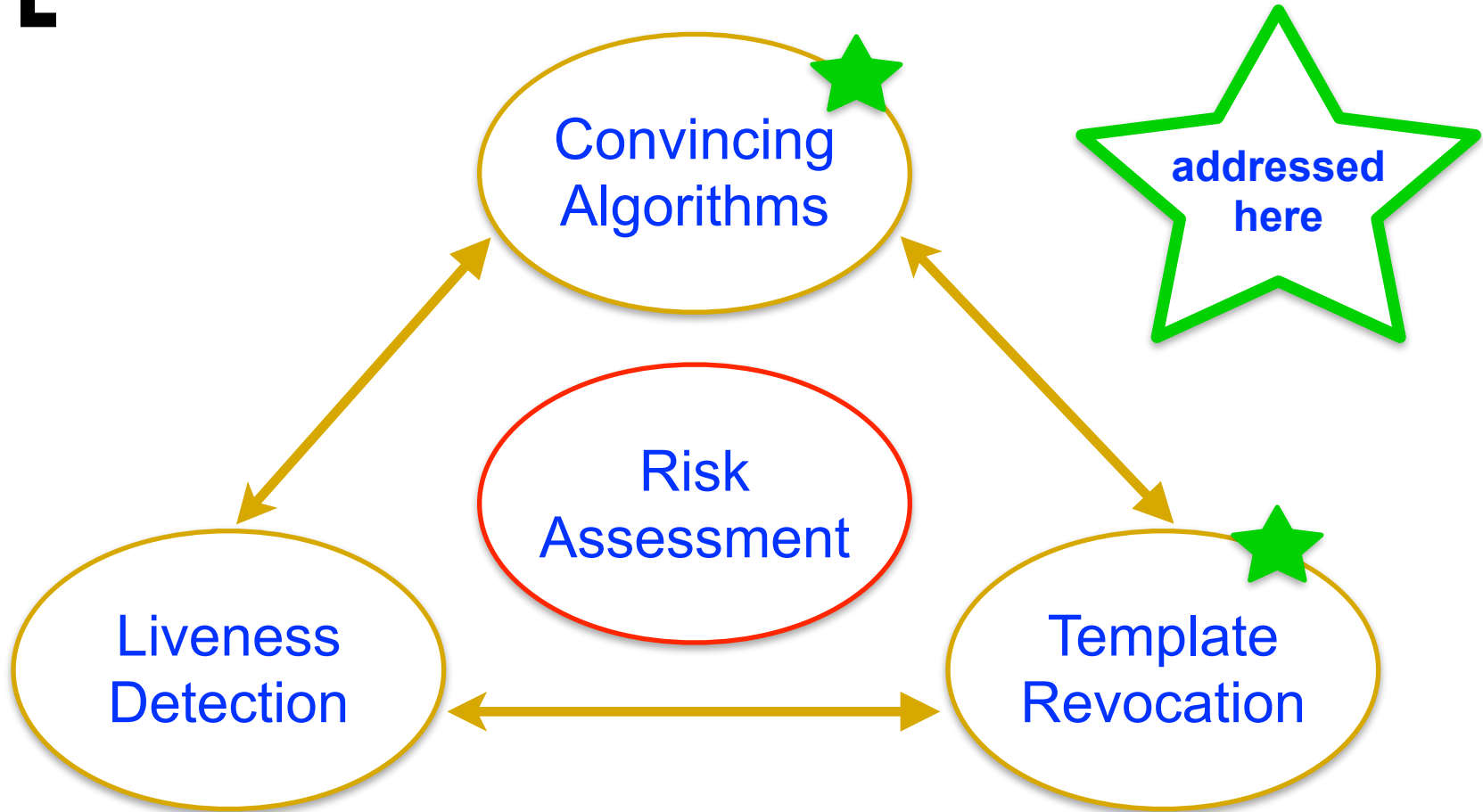
- *The process of defeating a biometric system through the introduction of fake biometric samples.*
  - *(Schuckers, Adler et al., 2010)*
- Particular modus operandi on how to deploy the attacking data vectors.
  - Can be seen as being orthogonal to the aforementioned hill-climbing attacks.

# [ Voice Biometrics Spoofing ]

- Spoofing techniques are, however, not “just helpers” as they are interesting on their own:
  - Text-To-Speech Synthesis
  - Voice Conversion
  - Artificial Signals



# [ Open Problems ]



# [ Biometrics In Mobile App ]

- Let's say we want to enhance a mobile banking application by biometrics.
- ...three-factor authentication by:
  - I) something to have (device key)
  - II) something to know (PIN)
  - III) something to be (BIO sample)

# Reflecting Privacy Protection



## Úřad pro ochranu osobních údajů

Pplk. Sochora 27, 170 00 Praha 7, Tel.: 234 665 111, Fax: 234 665 444; e-mail: posta@uouu.cz

### **STANOVISKO č. 3/2009**

květen 2009

#### **Biometrická identifikace nebo autentizace zaměstnanců**

##### **Úvod**

Záměrem stanoviska je vyjádřit základní přístupy Úřadu pro ochranu osobních údajů (dále jen „Úřad“) pro použití systémů umožňujících spolehlivé určení fyzické osoby na základě unikátních biometrických znaků, které se v poslední době velmi rozšířilo i v pracovněprávních vztazích. Nejčastěji je ze strany zaměstnavatele vznášen požadavek na poskytnutí otisků prstů (případně otisku dlaně) zaměstnanců pro použití v přístupových a docházkových systémech. Použití biometrických znaků má vyloučit možnosti klamání zaměstnavatele při použití jiných prostředků, např. identifikačních karet.

# [ Privacy Protection Conclusion ]

- There is a strong preference of biometric systems such that:
  - they do not process biometric samples left unintentionally
  - they do not store biometric template in one central database



# [ Local Templates ]

- We want to process the biometric data strictly locally in the mobile device.
  - So the bank does not store the precious BIO templates of its clients.
- Furthermore, we want to leverage the existing mechanism of distributed implicit PIN verification via (H)OTP.
  - cf. *“The Decline and Dawn of Two-Factor Authentication on Smart Phones”*, ISS 2012

# [ Naive Approach ]

```
sample = get_biometric_data();
```

```
if (match(sample, template) > eta)  
    continue_with_authentication();
```

```
else
```

```
    abort_authentication();
```

# [ Recall ATA ]

**Definition.** *Let the After-Theft Attack (ATA) be any attacking scenario that assumes the attacker has unlimited physical access to the user's smart device.*

- Imagine somebody steals your mobile phone...
- Despite being a really obvious threat, it is way too often neglected in contemporary applications.
- By a robbery, the attacker can even get access to unlocked screen or a paired computer, hence receiving another considerable favour!

# [ Naive Approach vs. ATA ]

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**bypassed!**

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    continue_with_authentication();
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```
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```

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```

**bypassed!**

**stolen!**

# [ Intermezzo ]

Recall how we process the PIN in mobile apps:

- i) unlock a  $PIN\_key$  by the PIN
  - ii) let  $MK = KDF(PIN\_key, device\_key)$
  - iii) verify  $MK$  with the bank using conventional crypto protocols
- ...distributed implicit PIN verification.***

# [ Intermezzo

PIN\_key is shared with the bank (not the PIN!)

Recall how we process the PIN in mobile apps:

- i) unlock a *PIN\_key* by the PIN
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- ...distributed implicit PIN verification.***

# [ Adding the BIO Factor ]

Is there something like “*BIO\_key*”?

We would have:

- i) unlock the *PIN\_key* by the PIN
- ii) unlock the *BIO\_key* by the user’s BIO
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- iv) verify *MK* with the bank using conventional crypto protocols

Again, *BIO\_key* is shared with the bank, not a BIO template

# [ Cryptography Exactness ]

Let  $y = AES_K(x)$  for a random  $K$ .

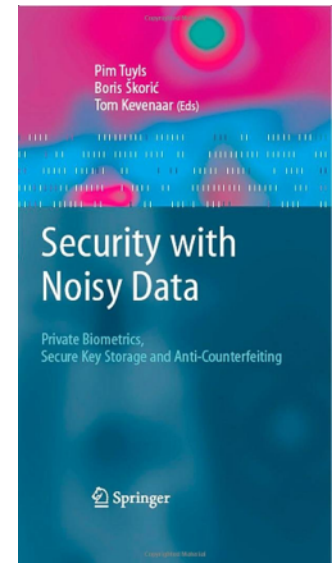
Then  $AES_K^{-1}(y) = x$ , while

$AES_{K \oplus 1}^{-1}(y) \neq x$  (probability  $\approx 1$ ).

- The better the algorithm is the more randomized response we get for even one-bit error.

# [ Biometrics Fuzziness ]

- We seldom get the same data in the subsequent scans of the very same person.
  - Actually, this is usually a clear sign of a spoofed sample.
- To overcome this (intra-user) variability, we can employ the *biometric cryptography*.





# [ BIO Cryptography ]

- Well, in 90's, there was a lot of alchemy in there.
  - Same as in crypto before C. E. Shannon in 1948 - 1949.
- Nowadays, it works hard towards a respected science.
- ...or how to deal with noisy data in cryptographic transformations.
  - These ideas go beyond the scope of biometrics. Quantum crypto or PUFs are further examples...
  - We can see the biometric cryptography as combining both feature quantization and classification into one “convolved” protocol.

# [ Our Illustrative Approach ]

- We employ BIO cryptography to cope with ATA threat in the mobile app.
- On behalf of this, we discuss the key concepts of these algorithms and protocols.

# [ Error-Correcting Code C ]

Let  $(F, \rho)$  be a metric space,  $\rho : F \times F \rightarrow [0, \infty)$ .

translation invariant metric:  $\rho(x, y) = \rho(0, x - y)$

Error correcting code is  $C \subset F, C = \{c_1, c_2, \dots\}$ .

*decode* :  $F \rightarrow C$

*t*-error correction capability:

Let  $\rho(c_i, y) \leq t$ , then *decode*( $c_i$ ) = *decode*( $y$ ) =  $c_i$ .

We assume *decode*() always returns a (possibly wrong) codeword.

# [ Metric For the Biometrics ]

- Let the extracted biometric features be expressible as an element of  $(F, \rho)$ .
  - Let also the  $\rho$ -distance measures the (dis)similarity of the two BIO samples.
    - We follow the *Fuzzy Commitment* by Juels and Wattenberg scheme that is a very good teaching example, since 1999.
    - It was (i.a.) generalised by Dodis et al. (2004) as *Fuzzy Extractor* based on *Secure Sketch*.
    - A well structured experiment exposing a particular ECC design to work with the iris code is by Hao et al. (2005).

# [ ECC Theory DO's and DON'Ts ]

- Recall, for ECC, we have solid proofs of guaranteed *random error correction capabilities*.
  - However, this is not the same as proofs of guaranteed *correlated error correction incapacibilities*.
- We need to combine low-level equation inspection together with overall statistics to get the assurance we want.

# [ Enrolment ]

- i) randomly choose  $c_{key} \in \mathbf{C} \subset \mathbf{F}$
- ii) get BIO features vector  $w \in \mathbf{F}$
- iii) let  $\xi = w - c_{key}$
- iv) let  $BIO\_key = hash(c_{key})$
- v) template =  $(\xi)$

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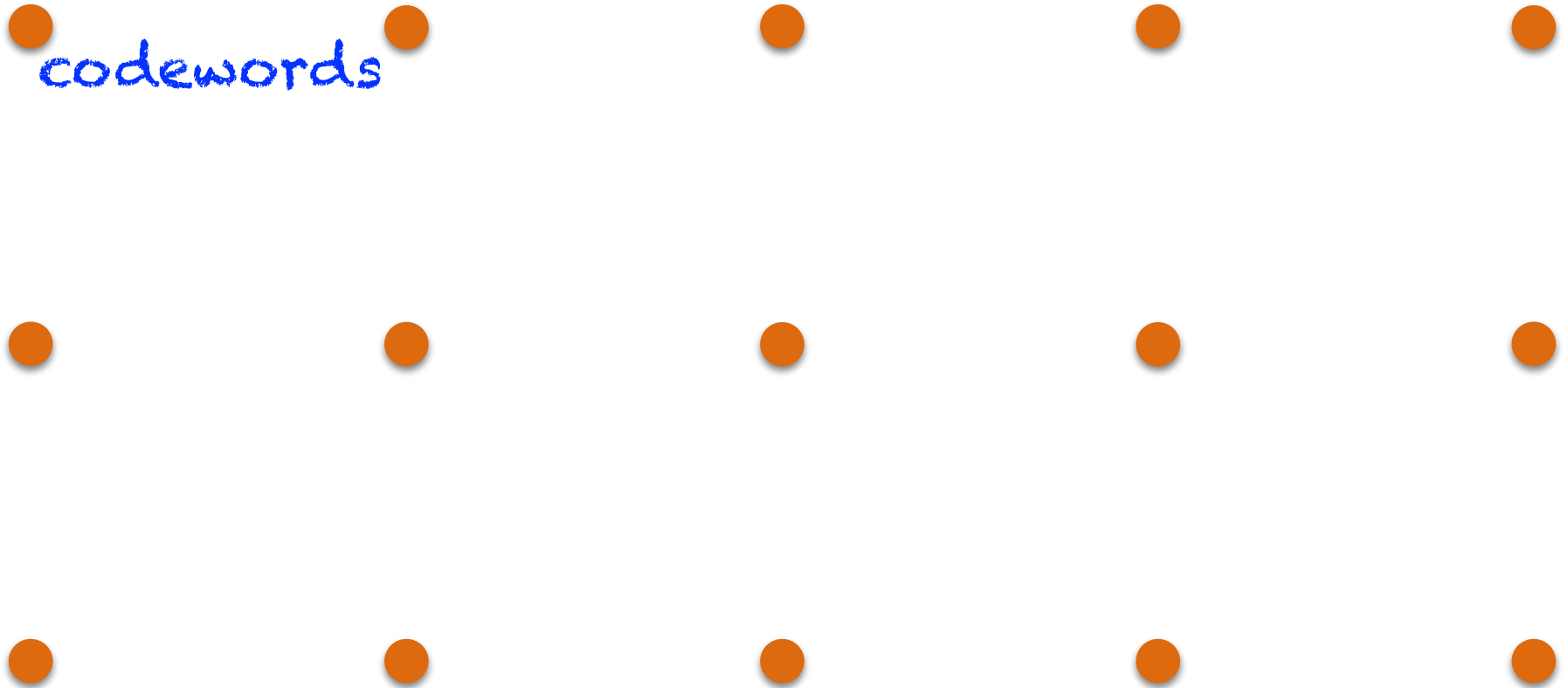
More involved entropy extractors can be used here...

# [ Verification ]

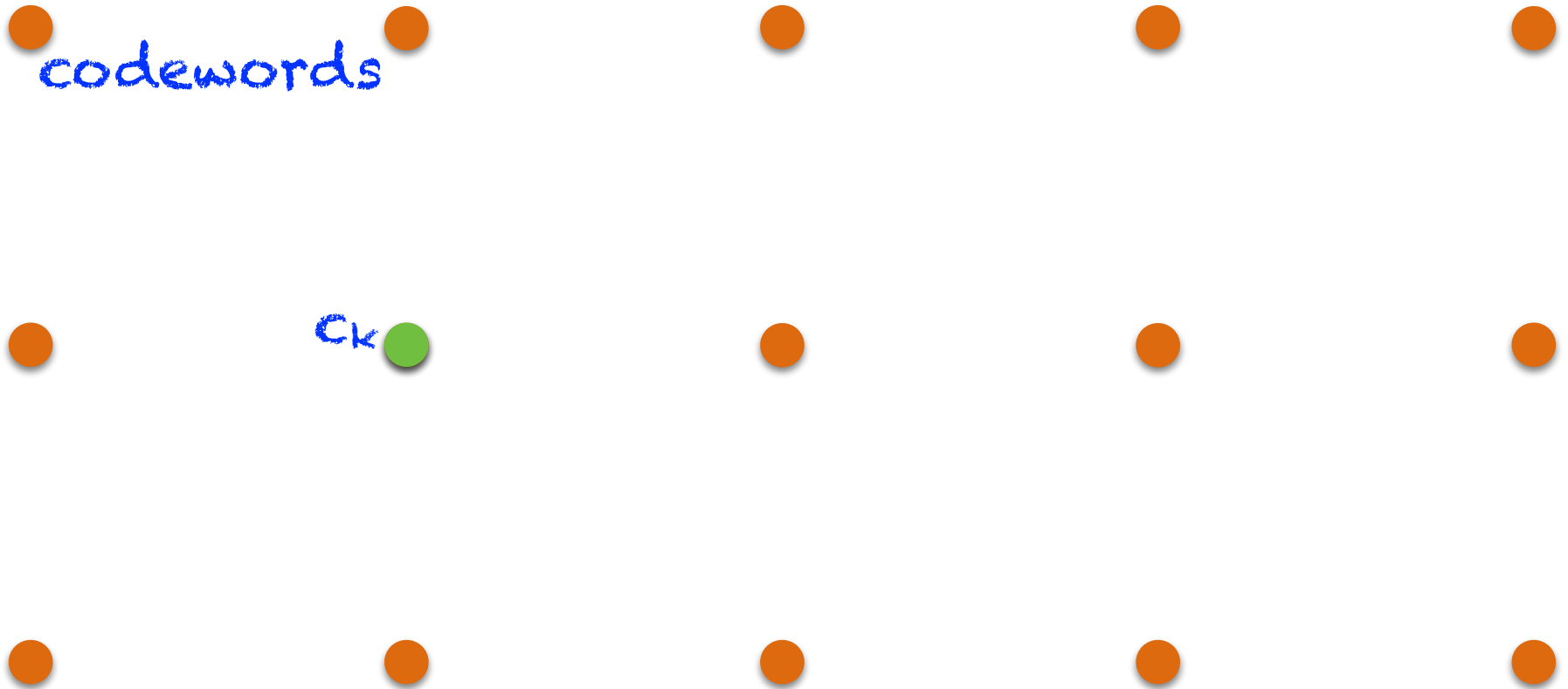
- i) get BIO features vector  $w' \in F$
- ii) let  $y = w' - \xi$
- iii) let  $c_{key}' = decode(y)$
- iv) let  $BIO\_key' = hash(c_{key}')$
- v) try to use  $BIO\_key'$  in the protocol above



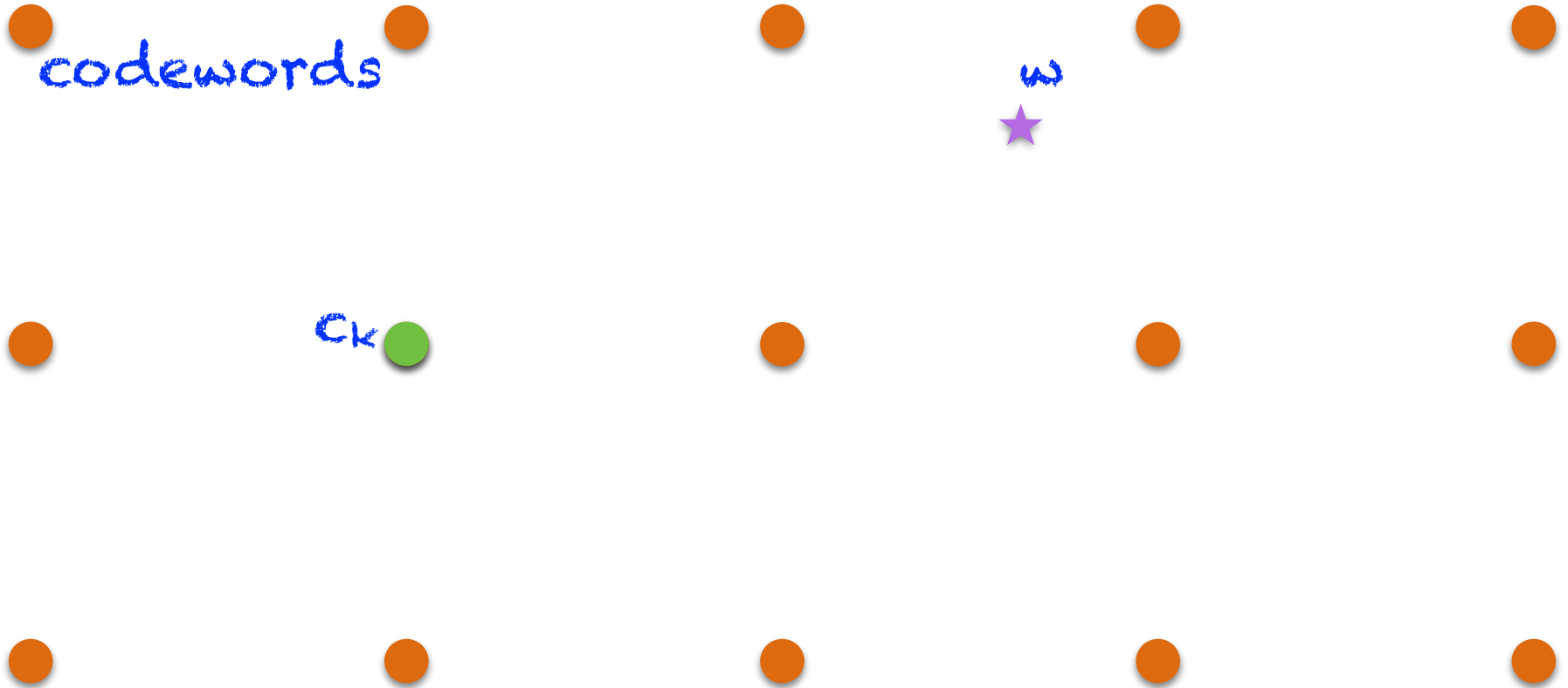
# [ Core Principle Illustrated ]



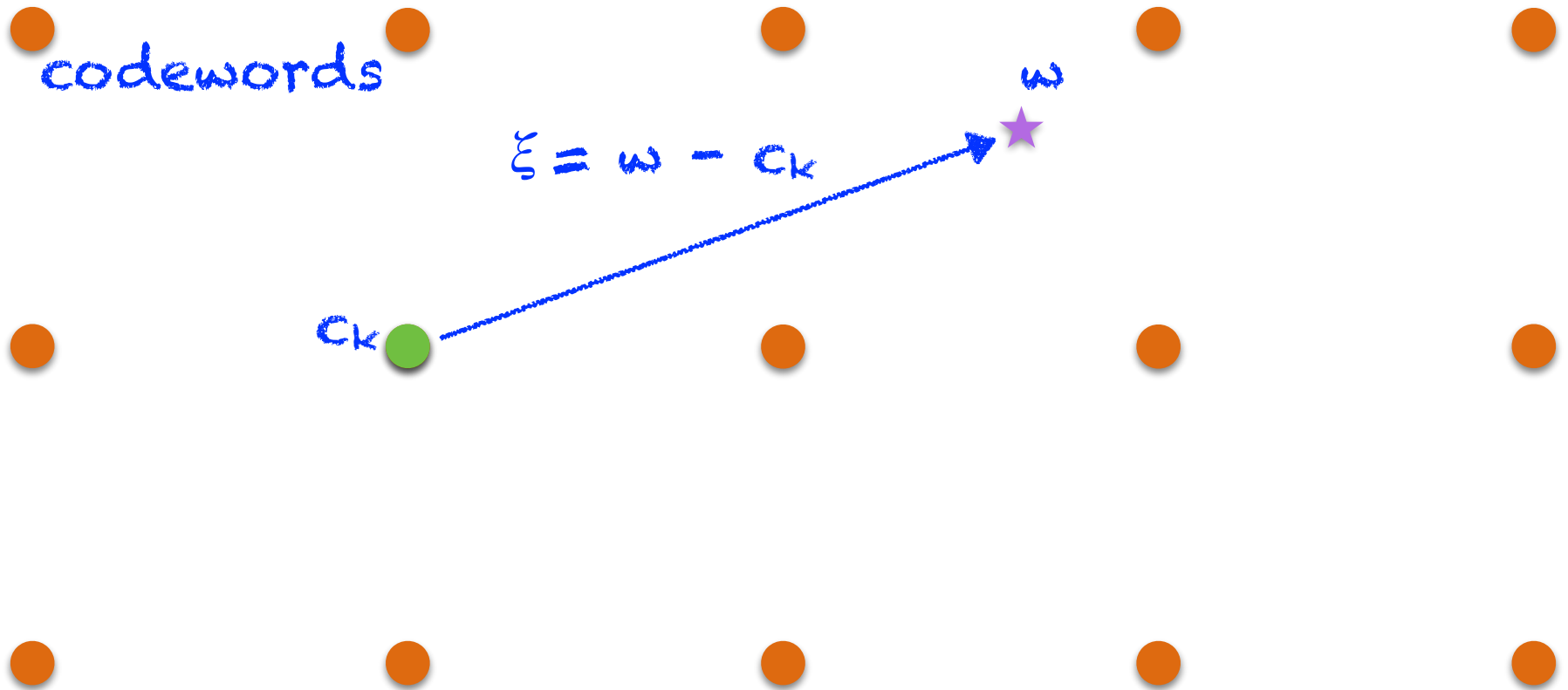
# [ Core Principle Illustrated ]



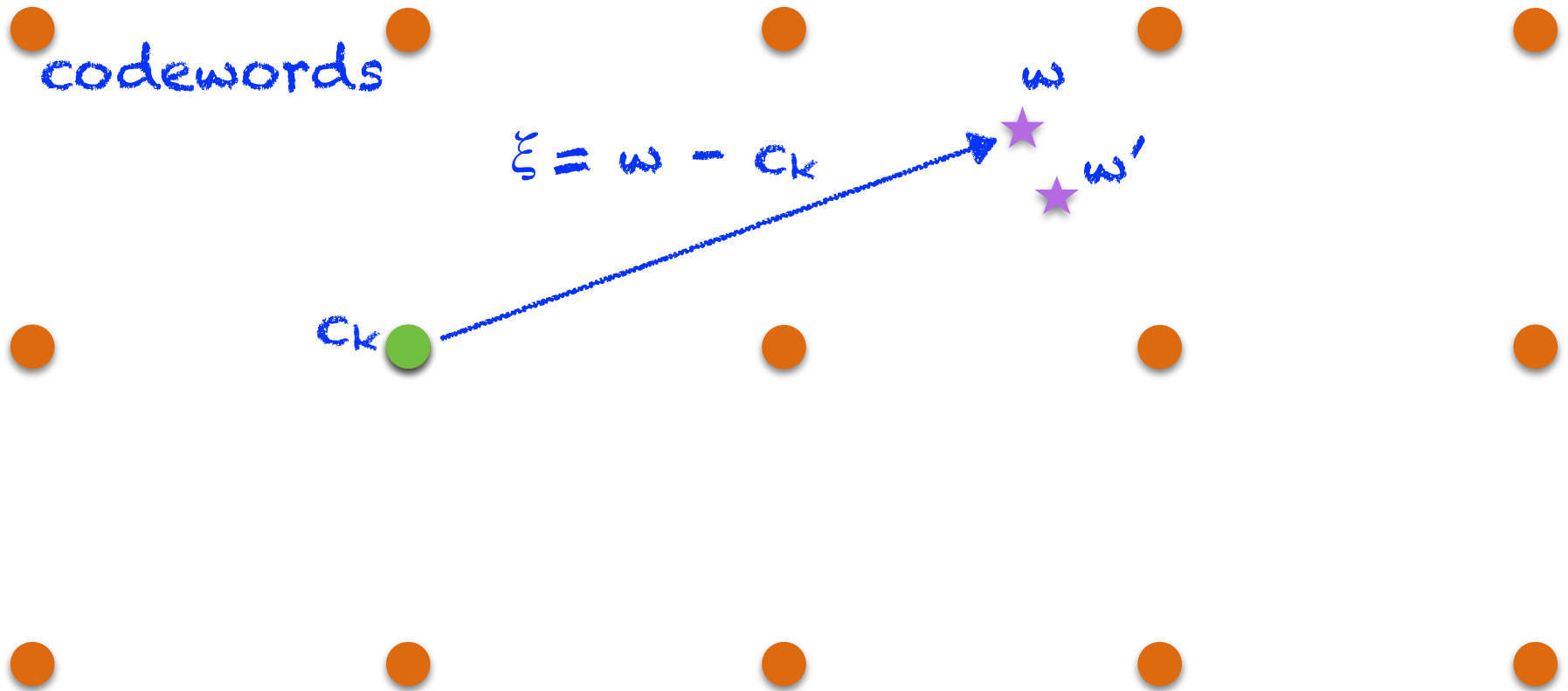
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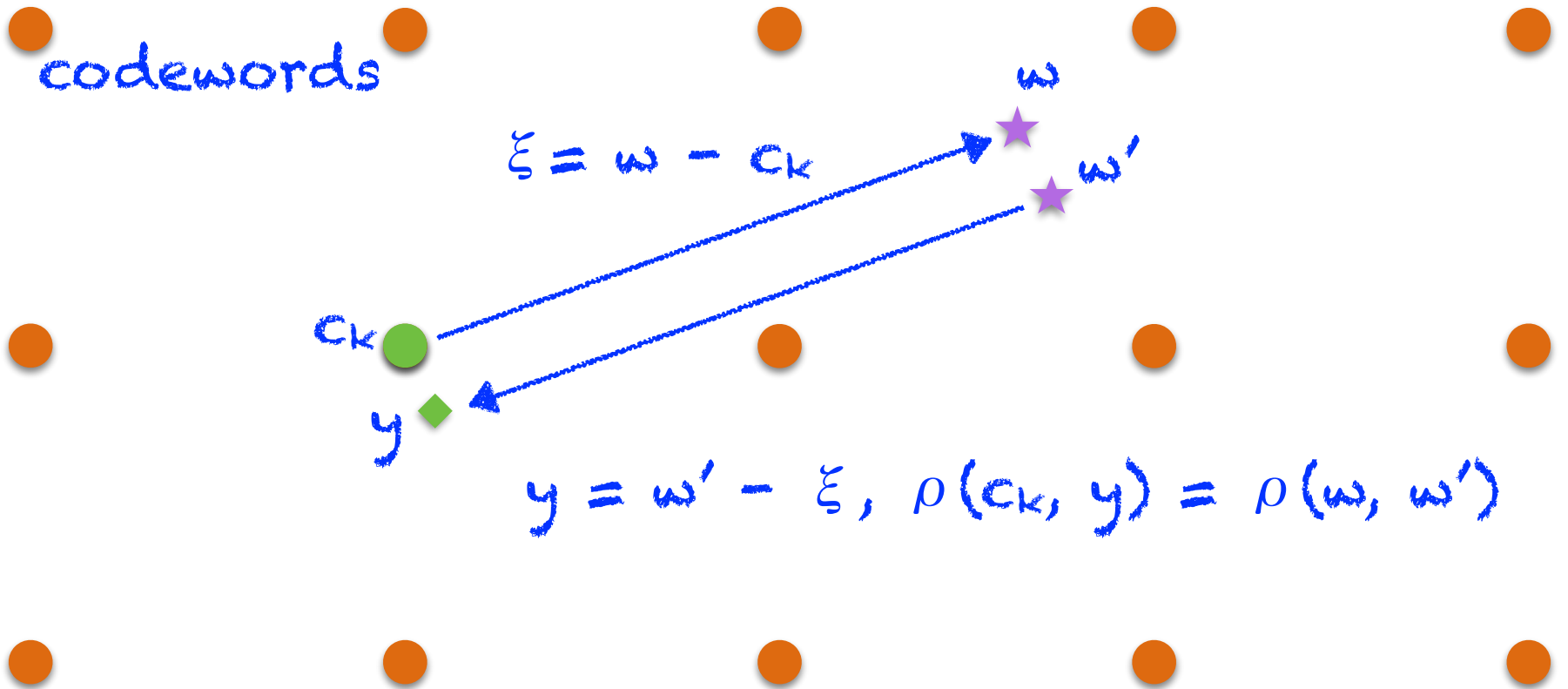
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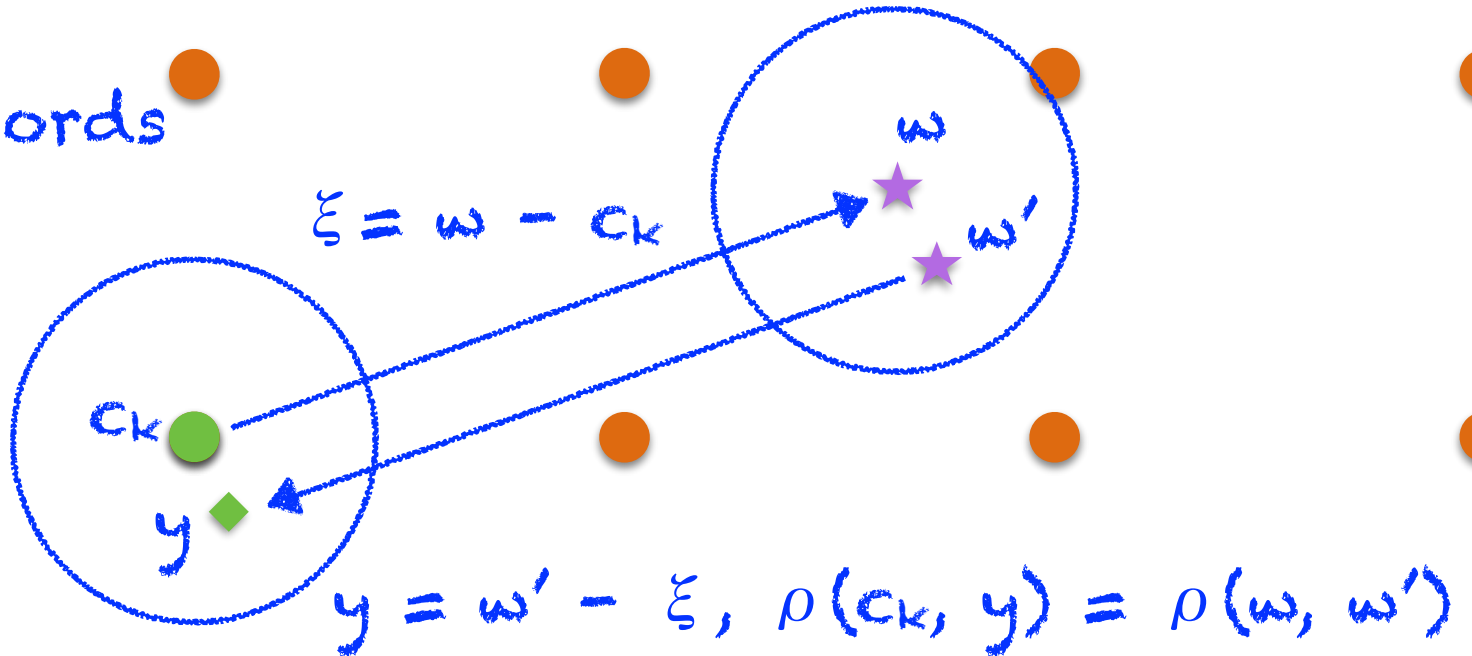


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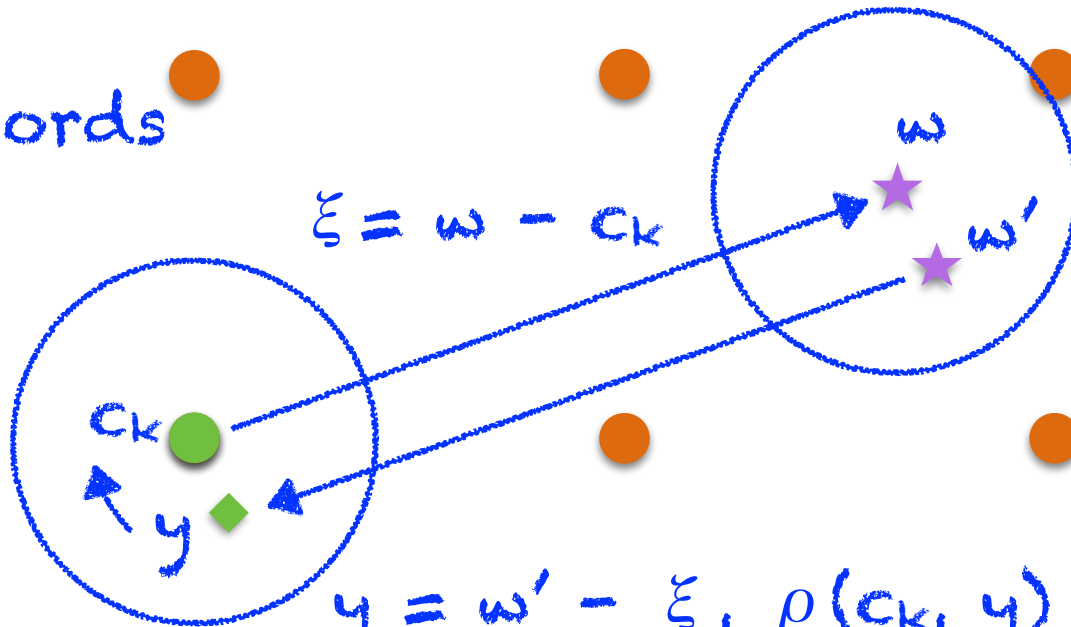
# [ Core Principle Illustrated ]

codewords



# [ Core Principle Illustrated ]

codewords



$$y = w' - \xi, \rho(c_k, y) = \rho(w, w')$$
$$\rho(w, w') \leq t \Rightarrow \text{decode}(y) = c_k$$



# [ Recovery Hint - $\xi$ ]

- Let  $D$  be the redundancy of the code  $C$  in  $F$  (with respect to randomly chosen codewords).
- Having learned  $\xi$ , the attacker gets at most  $D$  bits of information on the registration BIO sample  $w$ .
  - We emphasise, we do not store any hash-print of *BIO\_key* locally.
  - $\xi$  is the only information leaked under ATA.
  - Anyway, there are schemes allowing even local template encryption under a low-entropy password, cf. below.

# [ So, Is $\xi$ Public? ]

- Unless we have a plausible algebraic model for the biometric redundancy,  $\xi$  shall not be "public" as an RSA public key, for instance.
  - We rather suggest handling it the same way as the *device\_key* here.
  - Cf. also the encrypted template methods below.
- In our design, all the BIO cryptography is merely a life-saving jacket, not a silver bullet.
  - Yes, it is definitely important against ATA.
  - But we shall not overhype it!



My Voice Is My... *Entropy*

# [Voice-Based BIO-cryptography]

- We shall start with mapping the features of the whole utterance to a *supervector*  $w$ .
- We also have to enforce an ordering such that a particular coordinate of  $w$  always corresponds to a particular feature variable.
  - Straightforward for text-dependent methods.
  - For text-independent methods, we can follow the trick of Baum-Welch statistics re-ordering as employed in variants of Factor Analysis by Kenny, Dehak, Brümmer, et al.

# [Another BIO-Crypto Protocol]

- RBT ~ Randomized Biometric Templates
  - Ballard et al., 2008
  - Shares the basic idea of using an **error correction mechanism** to cope with intra-user variability.
  - Resulting RBT scheme can be viewed as a special kind of **Fuzzy Extractor**.
- Employs *randomized feature selection* together with plausible *template encryption* suitable for even a low-entropy password.

# [ RBT Password Protection ]

- The authors really strived hard to devise password-based protection of the whole RBT.
  - This way, the password entropy gets combined with the BIO entropy to considerably harden ATA.
- There shall be no *verifiable plaintexts* (Lomas et al. in 1989) in RBT, so we *could* use even our precious PIN here.
  - We shall, however, verify this with respect to the particular RBT calibration we would eventually use...

# [ Error Correction of RBT ]

- RBT employs a quantization of random variables for error correction.
  - This naturally introduces Euclidean distance metric for features variation.
- The role of the quantization boundary offset  $\alpha_i$  roughly corresponds to  $\xi$ .
  - Note that  $\alpha_i$  can be further transformed to a non-verifiable plaintext.
  - So, it can be protected by our precious PIN.

# [Voice-Based BKG]

- BKG ~ Biometric Key Generation
  - In 2010, Carrara and Adams described a voice-based BKG by using RBT and a novel extraction of *reliable features*.
  - Euclidean metric of RBT is highly welcome here.



# [Text Dependency]

- RBT assumes a strict order of the biometric features employed for the key derivation.
  - With the BKG based on *reliable features* extraction and RBT, this corresponds to the time order.
  - So, we get a text-dependent scheme.
- Using a feature vector derived by a variant of front-end **Factor Analysis**, we could, however, relax the time order to cover text-independent methods as well...

# [ Recall the Joint FA Model ]

$$M = m + \mathbf{U}x + \mathbf{V}y + \mathbf{D}z$$

# [ Recall the Joint FA Model ]

$$M = m + \mathbf{U}x + \mathbf{V}y + \mathbf{D}z$$

Speaker-specific features vector,  
we let  $w = y$ .

# [ Another Voice-Based Scheme ]

- In 2001-2002, Monroe et al. employed a strict quantization together with a **secret sharing scheme (SSS)** to:
  - cope with intra-speaker variation,
  - allow mixing the biometric randomness with a (possibly low-entropy) password.
    - this is done via template encryption while obeying the rule of no verifiable plaintexts

# [Text Dependency]

- To cope with ATA, the speech model part (besides the SSS) must be a speaker- and text-independent one.
  - But do not be fooled by this. This is merely to say there shall be no verifiable plaintexts (voiceprints).
  - The whole scheme, however, assumes the speaker is using the same utterance for both enrolment and key recovery.
  - ➔ Again, it is a text-dependent scheme.
  - ➔ Again, front-end **Factor Analysis** may provide us with a text-independent variant.

# Towards “Back-End” Order Invariance

- There is the **Fuzzy Vault** scheme by Juels and Sudan since 2002.
  - Instead of SSS, they employ a noisy polynomial reconstruction based on Reed-Solomon (de)coding.
  - Furthermore, they use the quantized features directly as  $x$ -coordinate “probes” for the secret polynomial.
  - Finally, they employ the idea of chaffing to conceal the correct  $(x, p(x))$  points.
- This scheme exhibits the important **order invariance** property, this time without front-end preprocessing tricks.
  - However, as for the VB the previous methods may be more appropriate even for TI schemes, despite the involved front-end preprocessing.

# [ Anyway, Fuzzy Extractors Take It All ]

- Dodis et al. shown **Fuzzy Vault** can be modelled and enhanced by the general **Fuzzy Extractor** approach (2004).
  - Their construction is based on the **set difference metric**.
  - It can be seen as an improved theoretical framework for the original FV construction.
  - The idea of using a noisy polynomial reconstruction stays the same.

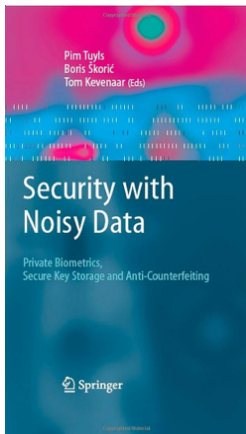
# [ Too Good To Be True? ]

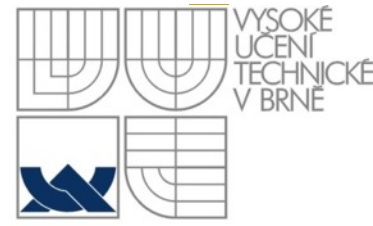
- The concise theory of **Security with Noisy Data** provides rather solid ground for robust protocols.
- We shall, however, verify the particular practical implementation very carefully.
  - There may be “surprisingly” new attacking strategies that were not incorporated in the former security “proofs” (Scheirer and Boulton, 2007) .
  - For instance, obtaining the recovery hints for multiple enrolments of the same individual may be a problem.
  - RBT cope with this by the random feature selection.
  - Distributed implicit *BIO\_key* verification also helps; **suitable entropy extractor** shall ensure *BIO\_key* is decorrelated from the original biometric data (to stop spreading it)!



# Conclusion

- **Fuzzy Extractors** together with the noisy data framework are the unifying theory of most of the BIO-cryptographic protocols.
  - The particular schemes developed more or less independently on FE then expose interesting practical tricks.
- To build up a real working system, we need to devise:
  - robust feature extraction,
  - error correction approach together with a suitable intra/inter variability metric,
  - key recovery and verification scheme,
  - template protection level (with a possible entropy boost from the client password/PIN).





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biocryptography, Brno, 2014

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